Structural Programming and Data Structures

Winter 2000

CMPUT 102: Sorting

Dr. Osmar R. Zaïane



University of Alberta

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Course Content

- Introduction
- Objects
- Methods
- Tracing Programs
- Object State
- Sharing resources
- Selection
- Repetition

- Vectors
- Testing/Debugging
- Arrays
- Searching
- Files I/O
- Sorting
- Inheritance
- Recursion



© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Objectives of Lecture 23Sorting

- Introduce the problem of sorting collections;
- Learn how to sort using a bubble sort algorithm;
- Learn how to sort with the selection algorithm.

Outline of Lecture 23



- The sorting problem
- Simple methods like bubble sort
- Selection sort example
- Selection sort code
- Complexity of selection sort



The Sort Problem

• Given a container, with elements that can be compared, put it in increasing or decreasing order.





| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 50 | 55 | 60 | 70 | 75 | 80 | 95 |

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Sorting Problem (con't)

- Given a container of *n* elements A[0..*n*-1] such that any elements x and y in the container A can be compared directly, either x<y, or x=y, or x>y.
- We want to permute the elements of A so that at the end A[0] ≤ A[1] ≤ ... ≤ A[n-1] (monotone non-decreasing), or A[0] ≥ A[1] ≥ ... ≥ A[n-1] (monotone decreasing)

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structur

University of Alberta



The Order of Things

Numbers

• Characters

$$\triangleright$$
 A < B < C < D < E < F < ... < X < Y < Z

$$\triangleright a < b < c < d < e < f < ... < x < y < z$$

$$\triangleright$$
a < z < A < Z

- Strings
 - > "Abacus" < "Alpha" < "Hello" < "Memorization" < "Memorize" < "Memory" < "Zebra"

Sorting

- There is often a need to put data in order.
- Sorting is among the most basic and universal of computational problems.
- There are hundreds of algorithms and variations on algorithms.
- Variety of sorting methods: internal vs. external, sorting in place vs. sorting with auxiliary structures, etc.

Outline of Lecture 23



- The sorting problem
- Simple methods like bubble sort
- Selection sort example
- Selection sort code
- Complexity of selection sort

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



One simple sorting method 35 18 22 97 61 10 Iterate over the collection and permute neighbours if necessary repeat iteration until no permutation possible. ↓ ↓ 18 35 22 97 61 10 ↓ ↓ 18 22 35 61 10 97 18 22 35 97 61 10 18 22 35 10 61 97 ↓ ↓ ... 18 22 35 97 61 10 ↓ ↓ ↓ ↓ ... 18 22 35 97 61 10 ↓ ↓ 10 18 22 35 61 97

© Dr. Osmar R. Zaïane, 2000

© Dr. Osmar R. Zaïane, 2000

18 22 35

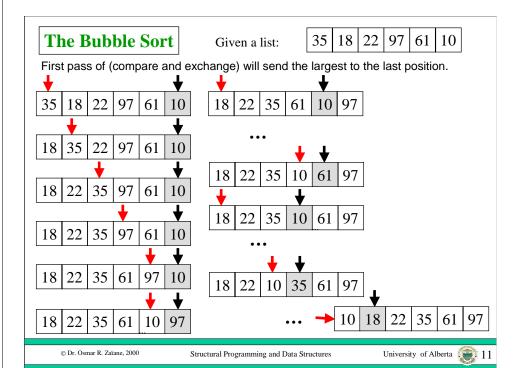
Structural Programming and Data Structures

University of Alberta

University of Alberta

10 | 18 | 22 | 35 | 61 | 97





```
public static void bubble_sort( int data[] ) {
    // Sort the given Array with selection sort method
    //(Ascending order)

    int current, last;

    for ( last = data.length-1; last >=1; last--)
        for ( current = 0; current < last; current++ )
        if ( data[current] > data[current+1] )
            this.exchange( data, current, current+1 );
    }
}
```

Structural Programming and Data Structures

Outline of Lecture 23



- The sorting problem
- Simple methods like bubble sort
- Selection sort example
- Selection sort code
- Complexity of selection sort

© Dr. Osmar R. Zaïane, 2000

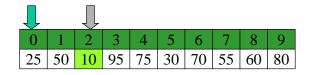
Structural Programming and Data Structures

University of Alberta



Selection Sort

• Look for the smallest element and exchange it with the element whose index is 0.



| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 50 | 25 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Selection Sort (con't)

• Look for the smallest element whose index is greater than or equal to 1 and exchange it with the element whose index is 1.

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 50 | 25 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

| | | | | | | | | | 9 |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 50 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures



Selection Sort (con't)

• Look for the smallest element whose index is greater than or equal to 2 and exchange it with the element whose index is 2.

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | | |
|----|----|----|----|----|----|----|----|----|----|--|--|
| 10 | 25 | 50 | 95 | 75 | 30 | 70 | 55 | 60 | 80 | | |
| | | | | | | | | | | | |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | | |
| 10 | 25 | 30 | 95 | 75 | 50 | 70 | 55 | 60 | 80 | | |

© Dr. Osmar R. Zaïane, 2000

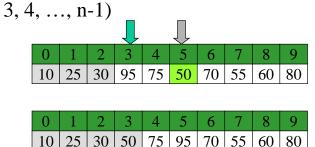
Structural Programming and Data Structur

University of Alberta



Selection Sort (con't)

• Look for the smallest element whose index is greater than or equal to k and exchange it with the element whose index is k (for k =



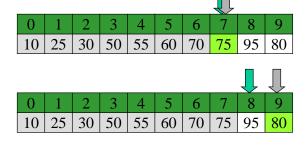
© Dr. Osmar R. Zaïane, 2000

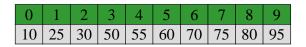
Structural Programming and Data Structures

University of Alberta



Selection Sort (con't)





Outline of Lecture 23



- The sorting problem
- Simple methods like bubble sort
- Selection sort example
- Selection sort code
- Complexity of selection sort



Selection Sort Algorithm

```
INPUT:
            data: an array of int
OUTPUT: data: sorted in ascending order
Method:
for ( first = 1; first < length - 1; first ++) {
   find Smallest such that data[Smallest] is the
   smallest between data[first] and data[length-1];
   permute Data[first] and Data[Smallest];
```

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Selection Sort Code

```
private void selectionSort(int anArray[]) {
// Sort the given Array with selection sort method (Ascending order)
               index;
  int
  int
              smallIndex;
  for (index = 0; index < anArray.length - 1; index++) {
       smallIndex = this.getSmallest(anArray, index);
       this.exchange(anArray, index, smallIndex);
```

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Code for method: exchange

```
private void exchange(int anArray[], int i, int j) {
// Exchange the elements of the array with
// the given two indexes.
   int
               temp;
  temp = anArray[i];
  anArray[i] = anArray[j];
  anArray[j] = temp;
```

Code for method: getSmallest

```
private int getSmallest(int anArray[], int start) {
// Return the index of the smallest element
// of the given array whose index is greater
// than or equal to the given start index.
               smallestIndex;
  int
  int
               index;
   smallestIndex = start:
  for (index = start + 1; index < anArray.length; index++)
       if (anArray[index] < anArray[smallestIndex])
               smallestIndex = index:
   return smallestIndex:
```

University of Alberta

Outline of Lecture 23



- The sorting problem
- Simple methods like bubble sort
- Selection sort example
- Selection sort code
- Complexity of selection sort

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta



Complexity of Selection Sort

- How many comparison operations are required for a selection sort of an *n*-element container?
- The sort method executes **getSmallest** for the indexes: 0, 1, ... n-2.
- Each time **getSmallest** is executed for an index, it does: (*n* - index) comparisons.
- The total number of comparisons is:

$$(n-0) + (n-1) + \dots + (n-(n-2)) = (1 + 2 + \dots + n) - 1 = n(n+1) - 1 \approx n^2$$
 for large n .

 $O(n^2)$ \rightarrow Quadratic time complexity

© Dr. Osmar R. Zaïane, 2000

Structural Programming and Data Structures

University of Alberta 26

